STAT253/317 Lecture 3: 4.3 Classification of States

Definition. Consider a Markov chain $\{X_n, n \ge 0\}$ with state space \mathfrak{X} . For two states $i, j \in \mathfrak{X}$, we say state j is accessible from state i if $P_{ii}^{(n)} > 0$ for some n, and we denote it as

$$i \rightarrow j$$
.

Note that **accessibility is transitive**: for $i, j, k \in \mathfrak{X}$, if $i \rightarrow j$ and $j \rightarrow k$, then $i \rightarrow k$.

Proof.

$$i \to j \quad \Rightarrow \quad P_{ij}^{(m)} > 0 \text{ for some } m$$

 $j \to k \quad \Rightarrow \quad P_{jk}^{(n)} > 0 \text{ for some } n$

By Chapman-Kolmogorov Equation:

$$P_{ik}^{(m+n)} = \sum_{l \in \mathfrak{X}} P_{il}^{(m)} P_{lk}^{(n)} \ge P_{ij}^{(m)} P_{jk}^{(n)} > 0,$$

which shows $i \rightarrow k$.

Communicability

Definition. Consider a Markov chain $\{X_n, n \ge 0\}$ chain with state space \mathfrak{X} . Two states $i, j \in \mathfrak{X}$ are said to **communicate** if $i \rightarrow j$, and $j \rightarrow i$. We denote it as

$$i \longleftrightarrow j$$
.

Fact. Communicability is also transitive, meaning that

if
$$i \longleftrightarrow j$$
 and $j \longleftrightarrow k$, then $i \longleftrightarrow k$.

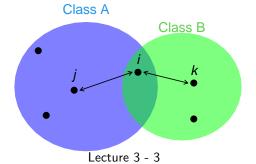
The proof is straight forward from the transitivity of accessibility.

Communicative Class

Definition. Two states that communicate with each other are in the same **class**. A state that communicates with no other states itself is a class.

Fact. Two classes are either identical or disjoint.

Proof. If two classes A and B have one state i in common, then all states in A communicate with i and all states in B do too. Consequently, all states with A can communicate with states in B (through state i). Class A and Class B must be identical.

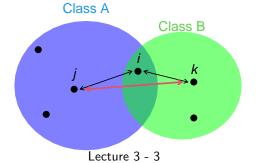


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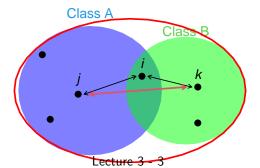


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Example 2. How many classes does the Ehrenfest diffusion model with K balls have?

$$\mathbb{P}_{1} = \begin{array}{cccccc} 1 & 2 & 3 & 4 & & 1 & 2 & 3 & 4 \\ 1 & 0.5 & 0.5 & 0 & 0 \\ 0.3 & 0.6 & 0.1 & 0 \\ 0 & 0 & 0.2 & 0.8 \\ 4 & 0 & 0 & 0.9 & 0.1 \end{array} \quad \mathbb{P}_{2} = \begin{array}{ccccccccc} 1 & 2 & 3 & 4 \\ 1 & 2 & 1/2 & 0 & 0 \\ 1/2 & 1/2 & 0 & 0 \\ 1/4 & 1/4 & 1/4 & 1/4 \\ 0 & 0 & 0 & 1 \end{array}$$

For \mathbb{P}_1 , $1 \leftrightarrow 2 \rightarrow 3 \leftrightarrow 4$. Classes: $\{1,2\}$, $\{3,4\}$.

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Example 2. How many classes does the Ehrenfest diffusion model with K balls have?

All states communicate. Only one class.

Closed Classes

Definition. A class C is said to be closed if

 $P_{ij} = 0$ for all *i* in *C* and *j* not in *C*.

Once the process gets into a closed class. It will never leave the class since the outgoing probabilities from the class are all 0.

Examples.

- For P₁ in the previous slide, the class {1,2} is not closed because it has a non-zero outgoing probability P₂3 = 0.1 > 0. The class {3,4} is closed.
- For P₂ in the previous slide, the classes {1, 2} and {4} are closed, and {3} is not closed.

A Markov Chain Restricted to a Closed Class is Also a Markov Chain

Example.

$$\mathbb{P}_1 = \begin{array}{cccc} 1 & 2 & 3 & 4 \\ 1 \\ 2 \\ 3 \\ 4 \end{array} \begin{pmatrix} 0.5 & 0.5 & 0 & 0 \\ 0.3 & 0.6 & 0.1 & 0 \\ 0 & 0 & 0.2 & 0.8 \\ 0 & 0 & 0.9 & 0.1 \end{pmatrix}$$

For P₁ above, the Markov chain restricted to the class {3,4} is also a Markov chain, with transition matrix

$$\begin{array}{rrrr}
3 & 4 \\
3 & \left(\begin{array}{cc}
0.2 & 0.8 \\
0.9 & 0.1
\end{array}\right)
\end{array}$$

► The Markov chain for P₁ can not be restricted to {1,2} as it may transit out of the state space from 2 to 3.

Irreducibility

A Markov chain is said to be **irreducible** if it has only 1 class.

Recurrence & Transience

Consider a Markov chain $\{X_n, n \ge 0\}$ chain with state space \mathfrak{X} . For $i \in \mathfrak{X}$, define

$$f_i = P(X_n = i \text{ for some } n > 0 | X_0 = i)$$

- If $f_i = 1$, we say state *i* is **recurrent** If $f_i < 1$, we say state *i* is **transient**
 - It's generally difficult to compute f_i directly.
 We need other tools to determine whether a state is recurrent or transient.

States in a Non-Closed Class Are Always Transient

For a class A that is NOT closed, there must exists some state k not in A such that

$$P_{i_0,k} > 0$$
, for some state i_0 in class A

but

$$P_{ki}^{(n)} = 0$$
 for all state *i* in class *A* and for all *n*.

Otherwise, state *i* would be accessible from state $k \ (k \longrightarrow i)$. As *i* and i_0 are in the same class, we know $i \leftrightarrow i_0$. Combining all the above, we have

$$k \longrightarrow i \longleftrightarrow i_0 \stackrel{P_{i_0,k}>0}{\longrightarrow} k.$$

and hence k would communicate with i_0 , contradicting to the assumption that k is not in A.

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Starting from a state j in a non-closed class A, there is a positive probability that the Markov chain will move to state k and never comes back to the class. Hence state j must be transient.

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Otherwise, state *i* would be accessible from state $k \pmod{i}$. As *i* and i_0 are in the same class, we know $i \leftrightarrow i_0$. Combining all the above, we have

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Are states in an closed class always recurrent?

Fact 1 If state *i* is recurrent, then starting from state *i*, the process will revisit state *i* infinitely often.

Fact 2 If state *i* is transient, then starting from state *i*, the number of times the process revisits state *i* is finite, with expected value $1/(1 - f_i)$. Reason: Let N_i be the number of times the process revisits state *i* after starting from *i*. Observe that

$$\begin{split} \mathrm{P}(N_i = k) &= \mathrm{P}(\text{returns to } i \text{ after 1st departure}) \\ &\times \cdots \times \mathrm{P}(\text{returns to } i \text{ after } k\text{th departure}) \\ &\times \mathrm{P}(\text{never returns to } i \text{ after } k\text{+}1\text{st departure}) \\ &= f_i^k(1 - f_i), \quad k = 0, 1, 2, \dots \end{split}$$

i.e., N_i has a geometric distribution with mean $1/(1 - f_i)$.

Claim:

$$\mathbb{E}(\# ext{ of visit to state } i|X_0=i) = \sum_{n=1}^\infty P_{ii}^{(n)}$$

Proof. Define

$$I_{ni} = \begin{cases} 1 & \text{if } X_n = i \\ 0 & \text{if } X_n \neq i \end{cases}, \quad n \ge 0, \ i \in \mathfrak{X}.$$

Observe that $\sum_{n=1}^{\infty} I_{ni}$ is the number of visits to state *i*.

$$\mathbb{E}\left[\sum_{n=1}^{\infty} I_{ni} \middle| X_0 = i\right] = \sum_{n=1}^{\infty} \mathbb{E}[I_{ni} \middle| X_0 = i]$$
$$= \sum_{n=1}^{\infty} P(X_n = i \middle| X_0 = i)$$
$$= \sum_{n=1}^{\infty} P_{ii}^{(n)}$$

Conclusion from Fact 1, Fact 2 and the Claim Above:

State *i* is recurrent $\iff \mathbb{E}(\# \text{ of visit to state } i | X_0 = i) = \infty.$

$$\iff \sum_{n=1}^{\infty} P_{ii}^{(n)} = \infty$$

Proposition 4.1

State *i* is
$$\begin{cases} \text{recurrent if} & \sum_{n=1}^{\infty} P_{ii}^{(n)} = \infty \\ \text{transient if} & \sum_{n=1}^{\infty} P_{ii}^{(n)} < \infty \end{cases}$$

Implication of Proposition 4.1:

States in a finite-state Markov chain CANNOT be all transient.

Reason. Observe that $\sum_{i \in X} I_{ni} = 1$ for all *n* since X_n must be in one of the states. Thus

$$\sum_{n=1}^{\infty}\sum_{i\in X}I_{ni}=\sum_{n=1}^{\infty}1=\infty.$$

Since \mathfrak{X} is finite, there exists at least one state *i* such that

$$\sum_{n=1}^{\infty}I_{ni}=\infty.$$

Such states are recurrent. Otherwise $\sum_{n=1}^{\infty} \sum_{i \in X} I_{ni}$ will be $< \infty$. Lecture 3 - 12

Corollary 4.2 If $i \leftrightarrow j$, and i is recurrent, then j is also recurrent. *Proof.*

$$i \to j \Rightarrow P_{ij}^{(k)} > 0$$
 for some k
 $j \to i \Rightarrow P_{ji}^{(l)} > 0$ for some l

By Chapman-Kolmogorov Equation:

$$\mathcal{P}_{jj}^{(l+n+k)} \geq \mathcal{P}_{ji}^{(l)} \mathcal{P}_{ii}^{(n)} \mathcal{P}_{ij}^{(k)}, ext{ for all } k=0,1,2,\ldots$$

Thus

$$\sum_{n=1}^{\infty} P_{jj}^{(n)} \ge \sum_{n=1}^{\infty} P_{jj}^{(l+n+k)} \ge \underbrace{P_{ji}^{(l)}}_{>0} \underbrace{\sum_{n=1}^{\infty} P_{ii}^{(n)}}_{=\infty} \underbrace{P_{ij}^{(k)}}_{>0} = \infty$$

Corollary 4.2 implies that all states of a finite irreducible Markov chain are recurrent.

Example: One-Dimensional Random Walk

$$X_{n+1} = egin{cases} X_n + 1 & ext{with prob. } p \ X_n - 1 & ext{with prob. } 1 - p \end{cases}$$

▶ State space
$$\{\cdots, -3, -2, -1, 0, 1, 2, 3, \cdots\}$$

All states communicate

$$\cdots \longleftrightarrow -2 \longleftrightarrow -1 \longleftrightarrow 0 \longleftrightarrow 1 \longleftrightarrow 2 \longleftrightarrow \cdots$$

It suffices to check whether 0 is recurrent or transient, i.e., whether $~~\sim$

$$\sum_{n=1}^{\infty}P_{00}^{(n)}=\infty ext{ or } <\infty$$

Example: One-Dimensional Random Walk (Cont'd)

$$P_{00}^{(2n+1)} = 0 \quad (Why?)$$

$$P_{00}^{(2n)} = {\binom{2n}{n}} p^n (1-p)^n$$

$$= \frac{(2n)!}{n! \, n!} p^n (1-p)^n \quad \text{Stirlin's Formula: } n! \approx n^{n+0.5} e^{-n} \sqrt{2\pi}$$

$$\approx \frac{(2n)^{2n+0.5} e^{-2n} \sqrt{2\pi}}{(n^{n+0.5} e^{-n} \sqrt{2\pi})^2} p^n (1-p)^n$$

$$= \frac{1}{\sqrt{\pi n}} [4p(1-p)]^n$$

Thus

$$\sum_{n=1}^{\infty} P_{ii}^{2n} \approx \sum_{n=1}^{\infty} \frac{1}{\sqrt{\pi n}} [4p(1-p)]^n \begin{cases} < \infty & \text{if } p \neq 1/2 \\ = \infty & \text{if } p = 1/2 \end{cases}$$

Conclusion: One-dimensional random walk is recurrent if p = 1/2, and transient otherwise.

Example: Two-Dimensional Symmetric Random Walk

Irreducible. Just check if 0 is recurrent.

$$P_{00}^{(2n)} = \sum_{i=0}^{n} \frac{(2n)!}{i!i!(n-i)!(n-i)!} \left(\frac{1}{4}\right)^{2n}$$
$$= \binom{2n}{n} \sum_{\substack{i=0\\n}}^{n} \binom{n}{i} \binom{n}{n-i} \left(\frac{1}{4}\right)^{2n}$$
$$= \binom{2n}{n}^{2} \left(\frac{1}{4}\right)^{2n} \approx \frac{1}{\pi n} \text{ by Stirlin's Formula}$$

Thus $\sum_{n=1}^{\infty} P_{00}^{(2n)} = \infty$.

Two-dimensional symmetric random walk is recurrent.

Example: *d*-Dimensional Symmetric Random Walk

In general, for a d-dimensional symmetric random walk, it can be shown that

$$P_{00}^{(2n)} \approx (1/2)^{d-1} \left(rac{d}{n\pi}
ight)^{d/2}$$

Thus

$$\sum_{n=1}^{\infty} P_{00}^{(2n)} \begin{cases} = \infty & \text{for } d = 1 \text{ or } 2 \\ < \infty & \text{for } d \ge 3 \end{cases}$$

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"A drunken man will find his way home. A drunken bird might be lost forever."