

GROTHENDIECK CONSTANT IS NORM OF STRASSEN MATRIX MULTIPLICATION TENSOR

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ABSTRACT. We show that two important quantities from two disparate areas of complexity theory — Strassen’s exponent of matrix multiplication ω and Grothendieck’s constant K_G — are intimately related. They are different measures of size for the same underlying object — the matrix multiplication tensor, i.e., the 3-tensor or bilinear operator $\mu_{l,m,n} : \mathbb{F}^{l \times m} \times \mathbb{F}^{m \times n} \rightarrow \mathbb{F}^{l \times n}$, $(A, B) \mapsto AB$ defined by matrix-matrix product over $\mathbb{F} = \mathbb{R}$ or \mathbb{C} . It is well-known that Strassen’s exponent of matrix multiplication is the greatest lower bound on (the log of) a *tensor rank* of $\mu_{l,m,n}$. We will show that Grothendieck’s constant is the least upper bound on a *tensor norm* of $\mu_{l,m,n}$, taken over all $l, m, n \in \mathbb{N}$. Aside from relating the two celebrated quantities, this insight allows us to rewrite Grothendieck’s inequality as a norm inequality

$$\|\mu_{l,m,n}\|_{1,2,\infty} = \max_{X,Y,M \neq 0} \frac{|\operatorname{tr}(XMY)|}{\|X\|_{1,2}\|Y\|_{2,\infty}\|M\|_{\infty,1}} \leq K_G.$$

We prove that Grothendieck’s inequality is unique: If we generalize the $(1, 2, \infty)$ -norm to arbitrary $p, q, r \in [1, \infty]$,

$$\|\mu_{l,m,n}\|_{p,q,r} = \max_{X,Y,M \neq 0} \frac{|\operatorname{tr}(XMY)|}{\|X\|_{p,q}\|Y\|_{q,r}\|M\|_{r,p}},$$

then $(p, q, r) = (1, 2, \infty)$ is, up to cyclic permutations, the only choice for which $\|\mu_{l,m,n}\|_{p,q,r}$ is uniformly bounded by a constant independent of l, m, n .

1. INTRODUCTION

Grothendieck’s inequality was originally established to relate fundamental norms on tensor product spaces [21]. Throughout this article, we will let $\mathbb{F} = \mathbb{R}$ or \mathbb{C} . The *Grothendieck constant* $K_G^{\mathbb{F}}$ is the sharp constant such that for every $l, m, n \in \mathbb{N}$ and every matrix $M = (M_{ij}) \in \mathbb{F}^{m \times n}$,

$$(1) \quad \max_{\|x_i\|=\|y_j\|=1} \left| \sum_{i=1}^m \sum_{j=1}^n M_{ij} \langle x_i, y_j \rangle \right| \leq K_G^{\mathbb{F}} \max_{|\varepsilon_i|=|\delta_j|=1} \left| \sum_{i=1}^m \sum_{j=1}^n M_{ij} \varepsilon_i \delta_j \right|$$

where the maximum on the left is taken over all $x_i, y_j \in \mathbb{F}^l$ of unit 2-norm, and the maximum on the right is taken over all $\varepsilon_i, \delta_j \in \mathbb{F}$ of unit absolute value (so over \mathbb{R} , $\varepsilon_i = \pm 1$ and $\delta_j = \pm 1$; over \mathbb{C} , $\varepsilon_i = e^{i\theta_i}$ and $\delta_j = e^{i\phi_j}$). The value on the left side of (1) is the same for all $l \geq m + n$ and as such some authors restrict themselves to $l = m + n$.

The existence of a such a constant independent of l, m and n was discovered by Alexandre Grothendieck in 1953. Alternative proofs via factorization of linear operators, geometry of Banach spaces, absolutely p -summing operators, etc, may be found in [40, 28, 38, 41] and references therein. In particular, the formulation in (1) was due to Lindenstrauss and Pełczyński [38].

The inequality has found applications in numerous areas, including Banach space theory, C^* algebra, harmonic analysis, operator theory, quantum mechanics, and most recently, computer science. In theoretical computer science, Grothendieck’s inequality has notably appeared in studies of

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unique games conjecture [29, 30, 31, 42, 43] and SDP relaxations of NP-hard combinatorial problems [2, 3, 4, 5, 11]. In quantum information theory, Grothendieck's inequality arises unexpectedly in Bell inequalities [17, 50, 24] and in XOR games [8, 9, 7], among several other areas; Grothendieck constants of specific orders, e.g., $K_G^{\mathbb{C}}(3)$ and $K_G^{\mathbb{C}}(4)$, also have important roles to play in quantum information theory [1, 25, 15]. The inequality has even been applied to some rather surprising areas, e.g., to communication complexity [39, 44, 45] and to privacy-preserving data analysis [16].

Although the Grothendieck constant appears in numerous mathematical statements and has many equivalent interpretations in physics and computer science, its exact value remains unknown and estimating increasingly sharper bounds for $K_G^{\mathbb{F}}$ has been a major undertaking. The current best known bounds are $K_G^{\mathbb{R}} \in [1.676, 1.782]$, established in [13] (lower) and [33] (upper); and $K_G^{\mathbb{C}} \in (1.338, 1.404]$, established in [14] (lower) and [22] (upper). A major recent breakthrough [6] established that Krivine's upper bound $\pi/(2 \log(1 + \sqrt{2})) \approx 1.782$ for $K_G^{\mathbb{R}}$ is not sharp. There have also been efforts in approximating Grothendieck's constants of specific orders, e.g., see [25, 15] for recent results on $K_G^{\mathbb{C}}(3)$ and $K_G^{\mathbb{C}}(4)$.

A world apart from the aforementioned areas touched by Grothendieck's inequality is the problem of complexity of matrix inversion, or equivalently, matrix multiplication, pioneered by Volker Strassen [48, 46, 49, 47]. A systematic study of this and other related problems has blossomed into what is now often called algebraic computational complexity [10]. For the uninitiated, Strassen famously discovered in [48] that the product of a pair of 2×2 matrices may be obtained with just seven multiplications:

$$\begin{bmatrix} a_1 & a_2 \\ a_3 & a_4 \end{bmatrix} \begin{bmatrix} b_1 & b_2 \\ b_3 & b_4 \end{bmatrix} = \begin{bmatrix} a_1 b_1 + a_2 b_2 & \beta + \gamma + (a_1 + a_2 - a_3 - a_4) b_4 \\ \alpha + \gamma + a_4 (b_2 + b_3 - b_1 - b_4) & \alpha + \beta + \gamma \end{bmatrix},$$

where $\alpha = (a_3 - a_1)(b_3 - b_4)$, $\beta = (a_3 + a_4)(b_3 - b_1)$, $\gamma = a_1 b_1 + (a_3 + a_4 - a_1)(b_1 + b_4 - b_3)$. Applied recursively, this gives an algorithm for forming the product of a pair of $n \times n$ matrices with just $O(n^{\log_2 7})$ multiplications, as opposed to $O(n^3)$ using the usual formula for matrix-matrix product. In addition, Strassen also showed that: (i) the number of additions may be bounded by a constant times the number of multiplications; (ii) matrix inversion may be achieved with the same complexity as matrix multiplication. In short, if there is an algorithm that forms matrix product in $O(n^\omega)$ multiplication, then it yields a $O(n^\omega)$ algorithm that would solve n linear equations in n unknowns, which is by far the most ubiquitous problem in all of scientific and engineering computing. The smallest possible ω became known as the *exponent of matrix multiplication*.

Strassen's astounding discovery captured the interests of numerical analysts and theoretical computer scientists alike and the complexity was gradually lowered over the years. Some milestones include the Coppersmith–Winograd [12] bound $O(n^{2.375477})$ that resisted progress for more than two decades until Vassilevska-Williams's improvement [51] to $O(n^{2.3728642})$; the current record, due to Le Gall [36], is $O(n^{2.3728639})$. Strassen showed [47] that the best possible ω is in fact given by

$$\omega = \inf_{n \in \mathbb{N}} \log_n(\text{rank}(\mu_{n,n,n})),$$

where $\mu_{n,n,n}$ is the *Strassen matrix multiplication tensor* — the 3-tensor in $(\mathbb{F}^{n \times n})^* \otimes (\mathbb{F}^{n \times n})^* \otimes \mathbb{F}^{n \times n}$ associated with matrix-matrix product, i.e., the bilinear operator

$$\mathbb{F}^{n \times n} \times \mathbb{F}^{n \times n} \rightarrow \mathbb{F}^{n \times n}, \quad (A, B) \mapsto AB.$$

Those unfamiliar with multilinear algebra [35] may regard the 3-tensor $\mu_{n,n,n}$ and the bilinear operator as the same object. If we choose a basis on $\mathbb{F}^{n \times n}$ (or three different bases, one on each copy of $\mathbb{F}^{n \times n}$), then $\mu_{n,n,n}$ may be represented as a 3-dimensional hypermatrix in $\mathbb{F}^{n^2 \times n^2 \times n^2}$. Over any \mathbb{F} -vector spaces $\mathbb{U}, \mathbb{V}, \mathbb{W}$, one may define *tensor rank* [26] for 3-tensors $\tau \in \mathbb{U} \otimes \mathbb{V} \otimes \mathbb{W}$ by

$$\text{rank}(\tau) = \min \left\{ r : \tau = \sum_{i=1}^r \lambda_i u_i \otimes v_i \otimes w_i \right\}.$$

In fact, Strassen showed that the tensor rank of a 3-tensor $\mu_\beta \in \mathbb{U}^* \otimes \mathbb{V}^* \otimes \mathbb{W}$ associated with a bilinear operator $\beta : \mathbb{U} \times \mathbb{V} \rightarrow \mathbb{W}$ gives the least number of multiplications required to compute β . The value of ω is in general dependent on the choice of \mathbb{F} , as tensor rank is well-known to be field dependent [37].

What exactly is ω ? The above discussion shows that it is the sharp lower bound for the (log of the) tensor rank of the Strassen matrix multiplication tensor:

$$(2) \quad \log_n(\text{rank}(\mu_{n,n,n})) \geq \omega \quad \text{for all } n \in \mathbb{N}.$$

What exactly is $K_G^{\mathbb{F}}$? We will show that it is the sharp upper bound for the tensor $(1, 2, \infty)$ -norm of the Strassen matrix multiplication tensor:

$$(3) \quad \|\mu_{l,m,n}\|_{1,2,\infty} \leq K_G^{\mathbb{F}} \quad \text{for all } l, m, n \in \mathbb{N}.$$

If we desire a greater parallel to (2), we may drop l and m in (3) — there is no loss of generality in assuming that $l = 2n$ and $m = n$, i.e., $K_G^{\mathbb{F}}$ is also the sharp upper bound so that

$$\|\mu_{2n,n,n}\|_{1,2,\infty} \leq K_G^{\mathbb{F}} \quad \text{for all } n \in \mathbb{N}.$$

In addition, the *Grothendieck constant of order* $l \in \mathbb{N}$, a popular notion in quantum information theory (e.g., [1, 15, 25]), is given by a simple variation, namely, the sharp upper bound $K_G^{\mathbb{F}}(l)$ in

$$\|\mu_{l,m,n}\|_{1,2,\infty} \leq K_G^{\mathbb{F}}(l) \quad \text{for all } m, n \in \mathbb{N}.$$

We will define the $(1, 2, \infty)$ -norm for an arbitrary 3-tensor formally in Section 4 but at this point it suffices to know its value for $\mu_{l,m,n}$, namely,

$$\|\mu_{l,m,n}\|_{1,2,\infty} = \max_{X,Y,M \neq 0} \frac{|\text{tr}(XMY)|}{\|X\|_{1,2}\|Y\|_{2,\infty}\|M\|_{\infty,1}}$$

where $X \in \mathbb{F}^{l \times m}$, $M \in \mathbb{F}^{m \times n}$, $Y \in \mathbb{F}^{n \times l}$, and $\|M\|_{p,q} := \max_{x \neq 0} \|Mx\|_q / \|x\|_p$ denotes the matrix (p, q) -norm.

The inequality (3) is in fact just Grothendieck's inequality. The characterizations of ω and K_G in (2) and (3) hold over both \mathbb{R} and \mathbb{C} although their values are field dependent. Incidentally the fact that Grothendieck's constant is essentially a tensor norm immediately explains why it is field dependent — because, as is the case for tensor rank, tensor norms are also field dependent [19].

An advantage of the formulation in (3) is that we obtain a natural family of (p, q, r) -norms on $\mu_{l,m,n}$ given by

$$\|\mu_{l,m,n}\|_{p,q,r} := \max_{X,Y,M \neq 0} \frac{|\text{tr}(XMY)|}{\|X\|_{p,q}\|Y\|_{q,r}\|M\|_{r,p}}$$

for any triple $1 \leq p, q, r \leq \infty$. This family of norms will serve as a platform for us to better comprehend Grothendieck's constant and Grothendieck's inequality. The study of the general (p, q, r) -case shows why the $(1, 2, \infty)$ -case is extraordinary. We will deduce a generalization of Grothendieck's inequality and show that the case $(p, q, r) = (1, 2, \infty)$, i.e., Grothendieck's inequality, is the only one up to trivial cyclic permutations¹ where there is a universal upper bound, i.e., Grothendieck's constant, that holds for all $l, m, n \in \mathbb{N}$.

Theorem 1.1 (Grothendieck–Hölder inequality). *Let $1 \leq p, q, r \leq \infty$ and $l, m, n \in \mathbb{N}$. Then*

$$\frac{1}{l^{1/q-1/2} \cdot m^{1/p-1/2} \cdot n^{1/r-1/2}} \leq \|\mu_{l,m,n}\|_{p,q,r} \leq K_G^{\mathbb{F}} \cdot l^{1/q-1/2} \cdot m^{1-1/p} \cdot n^{1/r}.$$

In particular, when $p = 1$, $q = 2$, and $r = \infty$, the upper bound gives Grothendieck's inequality (1).

Theorem 1.2 (Uniqueness of Grothendieck's inequality). *Let $1 \leq p, q, r \leq \infty$ and $l, m, n \in \mathbb{N}$. Then $\|\mu_{l,m,n}\|_{p,q,r}$ is uniformly bounded for all $l, m, n \in \mathbb{N}$ if and only if*

$$(p, q, r) \in \{(1, 2, \infty), (\infty, 1, 2), (2, \infty, 1)\}.$$

¹Unavoidable as (p, q, r) -norms are clearly invariant under cyclic permutations of p, q, r . See Lemma 4.1(i).

Theorem 1.1 follows from Theorems 4.2 and 4.3. Theorem 1.2 is just Theorem 5.2.

2. STRASSEN MATRIX MULTIPLICATION TENSOR

An important observation for us, obvious to anyone familiar with tensors [34, 35, 37] but perhaps less so to those accustomed to regarding (erroneously) a tensor as a “multiway array,” is that the bilinear operator

$$(4) \quad \beta \in \mathbb{F}^{l \times m} \times \mathbb{F}^{m \times n} \rightarrow \mathbb{F}^{l \times n}, \quad (X, Y) \mapsto XY,$$

and the trilinear functional

$$(5) \quad \tau : \mathbb{F}^{l \times m} \times \mathbb{F}^{m \times n} \times \mathbb{F}^{n \times l} \rightarrow \mathbb{F}, \quad (X, Y, Z) \mapsto \text{tr}(XYZ),$$

are given by² the same 3-tensor in

$$(\mathbb{F}^{l \times m})^* \otimes (\mathbb{F}^{m \times n})^* \otimes \mathbb{F}^{l \times n} \cong (\mathbb{F}^{l \times m})^* \otimes (\mathbb{F}^{m \times n})^* \otimes (\mathbb{F}^{n \times l})^*.$$

In other words, as 3-tensors, there is no difference between the product of two matrices and the trace of product of three matrices.

To see this, let $E_{ij} \in \mathbb{F}^{m \times n}$ denote the matrix with 1 in its (i, j) th entry and zeros everywhere else, so that $\{E_{ij} : i = 1, \dots, m; j = 1, \dots, n\}$ is the standard basis for $\mathbb{F}^{m \times n}$. Its dual basis for the dual space of linear functionals

$$(\mathbb{F}^{m \times n})^* := \{\varphi : \mathbb{F}^{m \times n} \rightarrow \mathbb{F} : \varphi(\alpha X + \beta Y) = \alpha\varphi(X) + \beta\varphi(Y)\}$$

is then given by $\{\varepsilon_{ij} : i = 1, \dots, m; j = 1, \dots, n\}$ where $\varepsilon_{ij} : \mathbb{F}^{m \times n} \rightarrow \mathbb{F}$, $X \mapsto x_{ij}$, is the linear functional that takes an $m \times n$ matrix to its (i, j) th entry. Now choose the standard inner product on $\mathbb{F}^{m \times n}$, i.e., $\langle X, Y \rangle = \text{tr}(X^T Y)$. Then $\varepsilon_{ij}(X) = \langle E_{ij}, X \rangle$ for all $X \in \mathbb{F}^{m \times n}$, which allows us to identify $(\mathbb{F}^{m \times n})^*$ with $\mathbb{F}^{n \times m}$ and linear functional $\varepsilon_{ij} \in (\mathbb{F}^{m \times n})^*$ with the matrix $E_{ji} \in \mathbb{F}^{n \times m}$.

It remains to observe that the usual formula for matrix-matrix product gives

$$\begin{aligned} \beta(X, Y) &= \sum_{i,k=1}^{l,n} \left(\sum_{j=1}^m x_{ij} y_{jk} \right) E_{ik} = \sum_{i,k=1}^{l,n} \left(\sum_{j=1}^m \varepsilon_{ij}(X) \varepsilon_{jk}(Y) \right) E_{ik} \\ &= \sum_{i,k=1}^{l,n} \left(\sum_{j=1}^m (\varepsilon_{ij} \otimes \varepsilon_{jk})(X, Y) \right) E_{ik} = \left(\sum_{i,j,k=1}^{l,m,n} \varepsilon_{ij} \otimes \varepsilon_{jk} \otimes E_{ik} \right) (X, Y), \end{aligned}$$

and thus

$$(6) \quad \beta = \sum_{i,j,k=1}^{l,m,n} \varepsilon_{ij} \otimes \varepsilon_{jk} \otimes E_{ik} \in (\mathbb{F}^{l \times m})^* \otimes (\mathbb{F}^{m \times n})^* \otimes \mathbb{F}^{l \times n}.$$

A similar simple calculation,

$$\begin{aligned} \tau(X, Y, Z) &= \sum_{i,j,k=1}^{l,m,n} x_{ij} y_{jk} z_{ki} = \sum_{i,j,k=1}^{l,m,n} \varepsilon_{ij}(X) \varepsilon_{jk}(Y) \varepsilon_{ki}(Z) \\ &= \sum_{i,j,k=1}^{l,m,n} (\varepsilon_{ij} \otimes \varepsilon_{jk} \otimes \varepsilon_{ki})(X, Y, Z) = \left(\sum_{i,j,k=1}^{l,m,n} \varepsilon_{ij} \otimes \varepsilon_{jk} \otimes \varepsilon_{ki} \right) (X, Y, Z), \end{aligned}$$

gives

$$(7) \quad \tau = \sum_{i,j,k=1}^{l,m,n} \varepsilon_{ij} \otimes \varepsilon_{jk} \otimes \varepsilon_{ki} \in (\mathbb{F}^{l \times m})^* \otimes (\mathbb{F}^{m \times n})^* \otimes (\mathbb{F}^{n \times l})^*.$$

By our identification, $(\mathbb{F}^{m \times n})^* = \mathbb{F}^{n \times m}$ and $\varepsilon_{ki} = E_{ik}$. So we see from (6) and (7) that indeed $\beta = \tau$ as 3-tensors. We denote this tensor by $\mu_{l,m,n}$. This has been variously called the *Strassen matrix multiplication tensor* or the *structure tensor for matrix-matrix product* [10, 34, 37, 52].

²To be more precise, by the universal property of tensor products [35, Chapter XVI, §1], β induces a linear map $\beta_* : \mathbb{F}^{l \times m} \otimes \mathbb{F}^{m \times n} \rightarrow \mathbb{F}^{l \times n}$ and τ induces a linear map $\tau_* : \mathbb{F}^{l \times m} \otimes \mathbb{F}^{m \times n} \otimes \mathbb{F}^{n \times l} \rightarrow \mathbb{F}$, i.e., $\beta_* \in (\mathbb{F}^{l \times m})^* \otimes (\mathbb{F}^{m \times n})^* \otimes \mathbb{F}^{l \times n}$ and $\tau_* \in (\mathbb{F}^{l \times m})^* \otimes (\mathbb{F}^{m \times n})^* \otimes (\mathbb{F}^{n \times l})^*$. We identify β, τ with the linear maps β_*, τ_* they induce.

3. GROTHENDIECK'S CONSTANT AND STRASSEN'S TENSOR

Let l, m, n be positive integers and let $M = (M_{ij}) \in \mathbb{F}^{m \times n}$. Let $x_1, \dots, x_m, y_1, \dots, y_n \in \mathbb{F}^l$ be vectors of unit 2-norm. We will regard x_1, \dots, x_m as columns of a matrix $X \in \mathbb{F}^{l \times m}$ and y_1^T, \dots, y_n^T as rows of a matrix $Y \in \mathbb{F}^{n \times l}$.

Recall that for any $p \geq 1$ with Hölder conjugate p^* , i.e., $1/p + 1/p^* = 1$, we have

$$(8) \quad \|X\|_{1,p} := \max_{z \neq 0} \frac{\|Xz\|_p}{\|z\|_1} = \max_{i=1, \dots, m} \|x_i\|_p, \quad \|Y\|_{p,\infty} := \max_{z \neq 0} \frac{\|Yz\|_\infty}{\|z\|_p} = \max_{i=1, \dots, n} \|y_i\|_{p^*},$$

and

$$\|M\|_{\infty,1} := \max_{z \neq 0} \frac{\|Mz\|_1}{\|z\|_\infty} = \max_{|\delta_j|=1} \sum_{i=1}^m \left| \sum_{j=1}^n M_{ij} \delta_j \right| = \max_{|\varepsilon_i|=1, |\delta_j|=1} \left| \sum_{i=1}^m \sum_{j=1}^n M_{ij} \varepsilon_i \delta_j \right|,$$

which may be further simplified for $\mathbb{F} = \mathbb{R}$ as

$$(9) \quad \|M\|_{\infty,1} = \max_{\varepsilon_i = \pm 1, \delta_j = \pm 1} \left| \sum_{i=1}^m \sum_{j=1}^n M_{ij} \varepsilon_i \delta_j \right| = \max_{\varepsilon, \delta \in \{\pm 1\}^n} |\varepsilon^T M \delta|.$$

We refer the reader to [19] for a proof that

$$(10) \quad \|\tau\|_{1,2,\infty} := \max_{X,Y,M \neq 0} \frac{|\tau(X, M, Y)|}{\|X\|_{1,2} \|Y\|_{2,\infty} \|M\|_{\infty,1}}$$

defines a norm for any tensor $\tau \in (\mathbb{F}^{l \times m})^* \otimes (\mathbb{F}^{m \times n})^* \otimes (\mathbb{F}^{n \times l})^*$, regarded as a trilinear functional.

Since

$$\sum_{i=1}^m \sum_{j=1}^n M_{ij} \langle x_i, y_j \rangle = \begin{cases} \text{tr}(XMY) & \text{if } \mathbb{F} = \mathbb{R}, \\ \text{tr}(XM\bar{Y}) & \text{if } \mathbb{F} = \mathbb{C}, \end{cases}$$

we see that Grothendieck's inequality (1) may be stated as

$$(11) \quad \max_{X,Y,M \neq 0} \frac{|\text{tr}(XMY)|}{\|X\|_{1,2} \|Y\|_{2,\infty} \|M\|_{\infty,1}} \leq K_G^{\mathbb{F}},$$

when $\mathbb{F} = \mathbb{R}$ and as

$$\max_{X,Y,M \neq 0} \frac{|\text{tr}(XM\bar{Y})|}{\|X\|_{1,2} \|Y\|_{2,\infty} \|M\|_{\infty,1}} \leq K_G^{\mathbb{F}},$$

when $\mathbb{F} = \mathbb{C}$. However, in the latter case, we may write

$$\max_{X,Y,M \neq 0} \frac{|\text{tr}(XM\bar{Y})|}{\|X\|_{1,2} \|Y\|_{2,\infty} \|M\|_{\infty,1}} = \max_{X,\bar{Y},M \neq 0} \frac{|\text{tr}(XM\bar{Y})|}{\|X\|_{1,2} \|\bar{Y}\|_{2,\infty} \|M\|_{\infty,1}} = \max_{X,Y,M \neq 0} \frac{|\text{tr}(XMY)|}{\|X\|_{1,2} \|Y\|_{2,\infty} \|M\|_{\infty,1}}$$

as matrix (p, q) -norms are invariant under complex conjugation. Hence (11) in fact gives Grothendieck's inequality for both $\mathbb{F} = \mathbb{R}$ and \mathbb{C} . By our discussion in Section 2 and our norm in (10), (11) is just

$$\|\mu_{l,m,n}\|_{1,2,\infty} \leq K_G^{\mathbb{F}}$$

where

$$(12) \quad \mu_{l,m,n} := \sum_{i=1}^l \sum_{j=1}^m \sum_{k=1}^n \varepsilon_{ij} \otimes \varepsilon_{jk} \otimes \varepsilon_{ki} \in (\mathbb{F}^{l \times m})^* \otimes (\mathbb{F}^{m \times n})^* \otimes (\mathbb{F}^{n \times l})^*$$

is the Strassen matrix multiplication tensor for the product of $l \times m$ and $m \times n$ matrices.

This allows us to define Grothendieck's constant in terms of tensor norms: For $\mathbb{F} = \mathbb{R}$ or \mathbb{C} ,

$$K_G^{\mathbb{F}} = \sup_{l,m,n \in \mathbb{N}} \|\mu_{l,m,n}\|_{1,2,\infty}.$$

Since $\|\mu_{l,m,n}\|_{1,2,\infty} = \|\mu_{m+n,m,n}\|_{1,2,\infty}$ for all $l \geq m+n$,

$$K_G^{\mathbb{F}} = \sup_{m,n \in \mathbb{N}} \|\mu_{m+n,m,n}\|_{1,2,\infty} = \sup_{n \in \mathbb{N}} \|\mu_{2n,n,n}\|_{1,2,\infty}.$$

In addition, the Grothendieck constant of order $l \in \mathbb{N}$ [1, 15, 25] may be defined as

$$K_G^{\mathbb{F}}(l) = \sup_{m,n \in \mathbb{N}} \|\mu_{l,m,n}\|_{1,2,\infty}.$$

4. GROTHENDIECK–HÖLDER INEQUALITY

The norm in (10) admits a natural generalization to arbitrary $p, q, r \in [1, \infty]$ as

$$\|\tau\|_{p,q,r} := \max_{X,Y,M \neq 0} \frac{|\tau(X, M, Y)|}{\|X\|_{p,q} \|Y\|_{q,r} \|M\|_{r,p}}$$

defined for any $\tau \in (\mathbb{F}^{l \times m})^* \otimes (\mathbb{F}^{m \times n})^* \otimes (\mathbb{F}^{n \times l})^*$, regarded as a trilinear functional

$$\tau : \mathbb{F}^{l \times m} \times \mathbb{F}^{m \times n} \times \mathbb{F}^{n \times l} \rightarrow \mathbb{F}.$$

In this article, we will only be interested in $\tau = \mu_{l,m,n}$, the Strassen tensor. We first state some simple observations that will be useful later.

Lemma 4.1. *Let $p, q, r \in [1, \infty]$. Then the (p, q, r) -norm of $\mu_{l,m,n}$*

(i) *is invariant under cyclic permutation of p, q, r ,*

$$\|\mu_{l,m,n}\|_{p,q,r} = \|\mu_{l,m,n}\|_{r,p,q} = \|\mu_{l,m,n}\|_{q,r,p};$$

(ii) *transforms under Hölder conjugation as*

$$\|\mu_{l,m,n}\|_{p,q,r} = \|\mu_{l,m,n}\|_{r^*,q^*,p^*}.$$

Recall that p^ is the Hölder conjugate of p , i.e., $1/p + 1/p^* = 1$.*

Proof. Since the numerator $\text{tr}(XMY) = \text{tr}(MYX) = \text{tr}(YXM)$ and the denominator is the product $\|X\|_{p,q} \|M\|_{r,p} \|Y\|_{q,r}$, cyclic permutations of (p, q) , (r, p) , (q, r) leave the quotient

$$\frac{|\text{tr}(XMY)|}{\|X\|_{p,q} \|M\|_{r,p} \|Y\|_{q,r}}$$

invariant. Now just observe that the cyclic permutations

$$(p, q), (r, p), (q, r) \rightarrow (q, r), (p, q), (r, p) \rightarrow (r, p), (q, r), (p, q)$$

correspond to the following permutations

$$(p, q, r) \rightarrow (q, r, p) \rightarrow (r, p, q).$$

Let X^\dagger denote the conjugate transpose of X . Since $|\text{tr}(XMY)| = |\text{tr}(Y^\dagger M^\dagger X^\dagger)|$ and $\|X\|_{p,q} = \|X^\dagger\|_{q^*,p^*}$, we have

$$\frac{|\text{tr}(XMY)|}{\|X\|_{p,q} \|Y\|_{q,r} \|M\|_{r,p}} = \frac{|\text{tr}(Y^\dagger M^\dagger X^\dagger)|}{\|Y^\dagger\|_{r^*,q^*} \|X^\dagger\|_{q^*,p^*} \|M^\dagger\|_{p^*,r^*}}.$$

Taking maximum over all nonzero X, Y, M yields the required equality. Note that the proof works over both \mathbb{R} and \mathbb{C} . \square

A straightforward application of Hölder's inequality yields an upper bound for $\|\mu_{l,m,n}\|_{p,q,r}$.

Theorem 4.2. *Let $p, q, r \in [1, \infty]$ and $l, m, n \in \mathbb{N}$. For any nonzero matrices $X \in \mathbb{F}^{l \times m}$, $Y \in \mathbb{F}^{n \times l}$ and $M \in \mathbb{F}^{m \times n}$, the following inequality is sharp:*

$$(13) \quad \frac{|\text{tr}(XMY)|}{\|X\|_{p,q} \|Y\|_{q,r} \|M\|_{r,p}} \leq \frac{|\text{tr}(XMY)|}{\|X\|_{1,2} \|Y\|_{2,\infty} \|M\|_{\infty,1}} \cdot l^{|1/q-1/2|} \cdot m^{1-1/p} \cdot n^{1/r}.$$

Furthermore, we have a generalization of Grothendieck's inequality:

$$(14) \quad \|\mu_{l,m,n}\|_{p,q,r} = \max_{X,Y,M \neq 0} \frac{|\text{tr}(XMY)|}{\|X\|_{p,q} \|Y\|_{q,r} \|M\|_{r,p}} \leq K_G^{\mathbb{F}} \cdot l^{|1/q-1/2|} \cdot m^{1-1/p} \cdot n^{1/r}.$$

Proof. First let $1 \leq q \leq 2$. Hölder's inequality together with the fact that $\|x\|_p \leq \|x\|_q$ whenever $q \leq p$ give us

$$(15) \quad \|X\|_{1,2} \leq \|X\|_{1,q} \leq \|X\|_{p,q} \quad \text{and} \quad \|Y\|_{2,\infty} \leq \|Y\|_{2,r} \leq l^{1/q-1/2} \|Y\|_{q,r}.$$

The same argument also gives $\|M\|_{\infty,p} \leq \|M\|_{\infty,1} \leq m^{1-1/p} \|M\|_{\infty,p}$ for $1 \leq p \leq \infty$ and thus

$$(16) \quad \|M\|_{\infty,1} \leq m^{1-1/p} \|M\|_{\infty,p} \leq n^{1/r} \cdot m^{1-1/p} \|M\|_{r,p}.$$

(13) then follows from (15) and (16). To see that it is sharp, we use the following³ $m \times n$ rank-one matrices:

$$E_{m,n} := \begin{bmatrix} 1 & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 \\ \vdots & \vdots & & \vdots \\ 0 & 0 & \dots & 0 \end{bmatrix}, \quad C_{m,n} := \begin{bmatrix} 1 & 0 & \dots & 0 \\ 1 & 0 & \dots & 0 \\ \vdots & \vdots & & \vdots \\ 1 & 0 & \dots & 0 \end{bmatrix}, \quad R_{m,n} := \begin{bmatrix} 1 & 1 & \dots & 1 \\ 0 & 0 & \dots & 0 \\ \vdots & \vdots & & \vdots \\ 0 & 0 & \dots & 0 \end{bmatrix}, \quad J_{m,n} := \begin{bmatrix} 1 & 1 & \dots & 1 \\ 1 & 1 & \dots & 1 \\ \vdots & \vdots & & \vdots \\ 1 & 1 & \dots & 1 \end{bmatrix}.$$

It is easy to check that

$$\begin{aligned} \|E_{l,m}\|_{p,q} &= 1, & \|R_{n,l}\|_{q,r} &= l^{1-1/q}, & \|J_{m,n}\|_{r,p} &= m^{1/p} \cdot n^{1-1/r}, \\ \|E_{l,m}\|_{1,2} &= 1, & \|R_{n,l}\|_{2,\infty} &= l^{1/2}, & \|J_{m,n}\|_{\infty,1} &= mn. \end{aligned}$$

Since (13) becomes an equality when $X = E_{l,m}$, $Y = R_{n,l}$, and $M = J_{m,n}$, it is sharp for $1 \leq q \leq 2$.

Next let $2 < q \leq \infty$. Similarly, we have

$$l^{1/q-1/2} \|X\|_{1,2} \leq \|X\|_{1,q} \leq \|X\|_{p,q} \quad \text{and} \quad \|Y\|_{2,\infty} \leq \|Y\|_{2,r} \leq \|Y\|_{q,r},$$

which together with (16) give us (13). In this case the sharpness follows from

$$\begin{aligned} \|C_{l,m}\|_{p,q} &= l^{1/q}, & \|E_{n,l}\|_{q,r} &= 1, & \|J_{m,n}\|_{r,p} &= m^{1/p} \cdot n^{1-1/r}, \\ \|C_{l,m}\|_{1,2} &= l^{1/2}, & \|E_{n,l}\|_{2,\infty} &= 1, & \|J_{m,n}\|_{\infty,1} &= mn, \end{aligned}$$

and selecting $X = C_{l,m}$, $Y = E_{n,l}$, and $M = J_{m,n}$.

(14) follows from taking maximum over nonzero X, M, Y and supremum over l, m, n . When $(p, q, r) = (1, 2, \infty)$, it yields Grothendieck's inequality (11). \square

The upper bound in (14) depends on l, m, n except when (p, q, r) is $(1, 2, \infty)$ or a cyclic permutation (by Lemma 4.1(i)). An immediate question is whether a uniform bound independent of l, m, n might perhaps also exist for some other values of (p, q, r) , i.e.,

$$(17) \quad K_{p,q,r} := \sup_{l,m,n \in \mathbb{N}} \|\mu_{l,m,n}\|_{p,q,r} < \infty?$$

In Section 5, we will see that $K_{p,q,r} = \infty$ for all $(p, q, r) \notin \{(1, 2, \infty), (\infty, 1, 2), (2, \infty, 1)\}$. Nevertheless, we stress that the absence of a uniform bound is only limited to the class of (p, q, r) -norms in (10). For example, we may consider the tensor *spectral norm* [19] of $\mu_{l,m,n}$,

$$\|\mu_{l,m,n}\|_{\sigma} := \max_{X,Y,M \neq 0} \frac{|\operatorname{tr}(XMY)|}{\|X\|_F \|Y\|_F \|M\|_F}$$

where the norm on X, Y, M is the matrix Frobenius (i.e., Hilbert–Schmidt) norm. In this case,

$$(18) \quad \|\mu_{l,m,n}\|_{\sigma} = 1, \quad \text{for all } l, m, n \in \mathbb{N},$$

since, by Cauchy–Schwartz and the submultiplicativity of the Frobenius norm,

$$|\operatorname{tr}(XMY)| \leq \|X\|_F \|MY\|_F \leq \|M\|_F \|X\|_F \|Y\|_F,$$

and equality is attained by choosing M, X, Y with 1 in the $(1, 1)$ th entry and 0 everywhere else.

We will use (18) to obtain lower bounds on $\|\mu_{l,m,n}\|_{p,q,r}$ below. (14) and (19) will collectively be referred to as the *Grothendieck–Hölder inequality*.

³These are standard in matrix theory, often used to demonstrate sharpness of various matrix inequalities.

Theorem 4.3. *Let $p, q, r \in [1, \infty]$ and $l, m, n \in \mathbb{N}$. Then*

$$(19) \quad \frac{1}{l^{|1/q-1/2|} \cdot m^{|1/p-1/2|} \cdot n^{|1/r-1/2|}} \leq \|\mu_{l,m,n}\|_{p,q,r}.$$

Proof. For $n \in \mathbb{N}$ and $p, q \in [1, \infty]$, let

$$c_{p,q}(n) := n^{\max\{0, 1/p-1/q\}}.$$

Then for any $M \in \mathbb{F}^{m \times n}$, the following sharp inequality holds [32, Theorem 4.3],

$$\|M\|_{p,q} \leq c_{q,2}(m)c_{2,p}(n)\|M\|_F.$$

It follows that

$$\|X\|_{p,q} \leq c_{q,2}(l)c_{2,p}(m)\|X\|_F, \quad \|Y\|_{q,r} \leq c_{r,2}(n)c_{2,q}(l)\|Y\|_F, \quad \|M\|_{r,p} \leq c_{p,2}(m)c_{2,r}(n)\|M\|_F,$$

and for any tensor $\tau \in (\mathbb{F}^{l \times m})^* \otimes (\mathbb{F}^{m \times n})^* \otimes (\mathbb{F}^{n \times l})^*$, we have

$$\|\tau\|_\sigma \leq \|\tau\|_{p,q,r} \cdot l^{|1/q-1/2|} \cdot m^{|1/p-1/2|} \cdot n^{|1/r-1/2|}.$$

Plugging in $\tau = \mu_{l,m,n}$ and using (18), we obtain (19). \square

A practical reason for wanting to ascertain (17) is that if

$$(20) \quad (p, q) \text{ and } (q, r) \in \{(1, 1), (2, 2), (\infty, \infty), (1, q), (q, \infty)\},$$

then $\|X\|_{p,q}$ and $\|Y\|_{q,r}$ can be computed in polynomial time (to arbitrary precision) and

$$\max_{X, Y \neq 0} \frac{|\text{tr}(XMY)|}{\|X\|_{p,q}\|Y\|_{q,r}} \leq K_{p,q,r}\|M\|_{r,p}$$

in principle gives a polynomial-time approximation of $\|M\|_{r,p}$, which is NP-hard [23] if (r, p) is not one of the special cases in (20). Unfortunately, we now know that as $K_{p,q,r} = \infty$ in all other cases, this only works when $(p, q, r) \in \{(1, 2, \infty), (\infty, 1, 2), (2, \infty, 1)\}$, all three of which are equivalent to Grothendieck's inequality.

5. GROTHENDIECK'S INEQUALITY IS UNIQUE

We show that $(p, q, r) = (1, 2, \infty)$ is, up to a cyclic permutation, the only case for which (17) holds. We will first rule out a large number of cases with the following proposition.

Proposition 5.1. *Let $p, q, r \in [1, \infty]$. If there exists a finite constant $K_{p,q,r} > 0$ such that $\|\mu_{l,m,n}\|_{p,q,r} \leq K_{p,q,r}$ for all $l, m, n \in \mathbb{N}$, then*

$$\min(p, q, r) = 1 \quad \text{and} \quad \max(p, q, r) = \infty.$$

Proof. Let $I_{m,n} \in \mathbb{F}^{m \times n}$ be the matrix obtained by appending zero rows or columns to the identity matrix⁴ I_n or I_m ,

$$I_{m,n} := \begin{cases} [I_n, 0_{m-n}]^\top & \text{if } m \geq n, \\ [I_m, 0_{n-m}] & \text{if } m < n. \end{cases}$$

Then its matrix (p, q) -norm is

$$(21) \quad \|I_{m,n}\|_{p,q} = \begin{cases} \min\{m, n\}^{1/q-1/p} & \text{if } p \geq q, \\ 1 & \text{if } p < q. \end{cases}$$

This follows from an easy calculation using Hölder inequality: For $m \geq n$,

$$\|I_{m,n}\|_{p,q} = \max_{z \neq 0} \frac{\|I_{m,n}z\|_q}{\|z\|_p} = \max_{z \neq 0} \frac{\|z\|_q}{\|z\|_p} = \begin{cases} n^{1/q-1/p} & \text{if } p \geq q, \\ 1 & \text{if } p < q, \end{cases}$$

⁴Note that $I_{n,n} = I_n$. For consistency, we will always use the latter notation when it is a square matrix.

and for $m < n$,

$$\|I_{m,n}\|_{p,q} = \max_{z \neq 0} \frac{\|I_{m,n}z\|_q}{\|z\|_p} = \max_{z \neq 0} \frac{\|z_m\|_q}{\|z\|_p} = \max_{z_m \neq 0} \frac{\|z_m\|_q}{\|z_m\|_p} = \begin{cases} m^{1/q-1/p} & \text{if } p \geq q, \\ 1 & \text{if } p < q, \end{cases}$$

where $z_m = [z_1, \dots, z_m] \in \mathbb{F}^m$ is the vector comprising the first m entries of z .

Set $X = I_{l,m}$, $Y = I_{n,l}$, and $M = I_{m,n}$. Then $\text{tr}(XMY) = \min\{l, m, n\}$, and by (21), we obtain

$$\frac{|\text{tr}(XMY)|}{\|X\|_{p,q}\|Y\|_{q,r}\|M\|_{r,p}} = \begin{cases} \min\{l, m, n\} \cdot \min\{m, n\}^{1/r-1/p} & \text{if } p \leq q \leq r, \\ \min\{l, m, n\} \cdot \min\{l, n\}^{1/q-1/r} \cdot \min\{m, n\}^{1/r-1/p} & \text{if } p \leq r \leq q, \\ \min\{l, m, n\} \cdot \min\{l, m\}^{1/p-1/q} \cdot \min\{m, n\}^{1/r-1/p} & \text{if } q \leq p \leq r, \\ \min\{l, m, n\} \cdot \min\{l, m\}^{1/p-1/q} & \text{if } q \leq r \leq p, \\ \min\{l, m, n\} \cdot \min\{l, n\}^{1/q-1/r} & \text{if } r \leq p \leq q, \\ \min\{l, m, n\} \cdot \min\{l, m\}^{1/p-1/q} \cdot \min\{l, n\}^{1/q-1/r} & \text{if } r \leq q \leq p. \end{cases}$$

Suppose $l = 2n$, $m = n$ and $p \leq q \leq r$, then

$$\lim_{n \rightarrow \infty} \frac{|\text{tr}(XMY)|}{\|X\|_{p,q}\|Y\|_{q,r}\|M\|_{r,p}} = \lim_{n \rightarrow \infty} n^{1/r-1/p+1} = \infty$$

unless $p = 1$ and $r = \infty$. Repeating the argument for all possible permutations of (p, q, r) and taking advantage of Lemma 4.1(i), we conclude that $\min(p, q, r) = 1$ and $\max(p, q, r) = \infty$ is necessary for the uniform boundedness of $\|\mu_{l,m,n}\|_{p,q,r}$. \square

We will next eliminate the remaining possibilities. Our approach will rely on the existence of Hadamard matrices of arbitrarily large dimensions. Indeed, an $n \times n$ Hadamard matrix $H_n \in \{\pm 1\}^{n \times n}$ exists for any n divisible by 4, or, for concreteness, we may set $H_n = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}^{\otimes k}$ with $n = 2^k$ [27, Section 2.1]. The bottom line is that we may let $n \rightarrow \infty$ in the proof below.

Theorem 5.2 (Uniqueness of Grothendieck's inequality). *Let $1 \leq p, q, r \leq \infty$ and $l, m, n \in \mathbb{N}$. Then $\|\mu_{l,m,n}\|_{p,q,r}$ is uniformly bounded for all $l, m, n \in \mathbb{N}$ if and only if*

$$(p, q, r) \in \{(1, 2, \infty), (\infty, 1, 2), (2, \infty, 1)\}.$$

Proof. We will see that it suffices to take $l = m = n$ throughout this proof. By Lemma 4.1(i) and Proposition 5.1, we may assume that $p = 1$ and either $q = \infty$ or $r = \infty$. We will show that $\text{tr}(XMY)$ is unbounded for judiciously chosen $n \times n$ real matrices X , M , and Y as $n \rightarrow \infty$.

CASE I: $(1, q, \infty)$, $1 \leq q \leq \infty$. Suppose $2 < q \leq \infty$. Let $X_0 = n^{-1/q}\Delta$ for some arbitrary $\Delta = (\delta_{ij}) \in \{\pm 1\}^{n \times n}$ and let $Y_0 = I_n$. Then $\|X_0\|_{1,q} = n^{-1/q}\|\Delta\|_{1,q} = 1$ and $\|Y_0\|_{q,\infty} = \|I_n\|_{q,\infty} = 1$ by (8). For any $M = (M_{ij}) \in \mathbb{R}^{n \times n}$,

$$\max_{\|X\|_{1,q}, \|Y\|_{q,\infty} \leq 1} |\text{tr}(XMY)| \geq |\text{tr}(X_0MY_0)| = n^{-1/q}|\text{tr}(\Delta M)| = n^{-1/q} \left| \sum_{i,j=1}^n \delta_{ij} M_{ij} \right|.$$

Since $\Delta \in \{\pm 1\}^{n \times n}$ is arbitrary, we will choose δ_{ij} so that $\delta_{ij}M_{ij}$ is nonnegative. Thus

$$(22) \quad \max_{\|X\|_{1,q}, \|Y\|_{q,\infty} \leq 1} |\text{tr}(XMY)| \geq n^{-1/q} \sum_{i,j=1}^n |M_{ij}|.$$

Let $H_n \in \{\pm 1\}^{n \times n}$ be a Hadamard matrix. So $H_n H_n^T = nI_n$ and all singular values of H_n are \sqrt{n} [18]. Therefore, by (9),

$$(23) \quad \|H_n\|_{\infty,1} = \max_{\varepsilon, \delta \in \{\pm 1\}^n} |\varepsilon^T H_n \delta| \leq \sigma_{\max}(H_n) \|\varepsilon\|_2 \|\delta\|_2 = n^{3/2}.$$

Let $M = n^{-3/2}H_n$. Then $\|M\|_{\infty,1} \leq 1$ and by (22),

$$\max_{\|X\|_{1,q}, \|Y\|_{q,\infty}, \|M\|_{\infty,1} \leq 1} |\text{tr}(XMY)| \geq n^{-1/q} \times n^{-3/2} \times n^2 = n^{1/2-1/q} \rightarrow \infty$$

as $n \rightarrow \infty$.

Suppose $1 \leq q < 2$. Since the Hölder conjugates are $r^* = 1$, $2 < q^* \leq \infty$, and $p^* = \infty$, by Lemma 4.1(ii), this reduces to the case we just treated.

CASE II: $(1, \infty, r)$, $1 \leq r \leq \infty$. For $r = \infty$, we have $(1, \infty, \infty)$, which is same as the $q = \infty$ case in CASE I. For $r = 1$, we have $(1, \infty, 1)$, but by Lemma 4.1(i), this is equivalent to $(1, 1, \infty)$, which is same as the $q = 1$ case in CASE I. So we may assume $1 < r < \infty$.

Suppose $1 < r < 2$. Let $M = n^{1/r-1}I_n$ and $Y = n^{-3/2}H_n$ where $H_n \in \{\pm 1\}^{n \times n}$ is a Hadamard matrix. Then $\|M\|_{r,1} = n^{1/r-1}\|I_n\|_{r,1} = 1$ by (21), and $\|Y\|_{\infty,r} \leq \|Y\|_{\infty,1} = n^{-3/2}\|H_n\|_{\infty,1} \leq 1$ by (23). We choose $X \in \{\pm 1\}^{n \times n}$ such that $\text{tr}(XH_n) = n^2$ and thus $\text{tr}(XY) = n^{-1/2}$. Clearly $\|X\|_{1,\infty} = 1$ by (8). Hence

$$\text{tr}(XMY) = n^{1/r-1} \text{tr}(XY) = n^{1/r-1/2} \rightarrow \infty$$

as $n \rightarrow \infty$.

Suppose $2 < r < \infty$. Since the Hölder conjugates are $1 < r^* < 2$, $q^* = 1$, and $p^* = \infty$, by Lemma 4.1(ii), this is equivalent to the case $(r^*, 1, \infty)$. Now by Lemma 4.1(i), this is in turn equivalent to the case $(1, \infty, r^*)$ with $1 < r^* < 2$, which is the case we just treated.

Suppose $r = 2$. Let $Y = n^{-1}H_n$ where $H_n \in \{\pm 1\}^{n \times n}$ is again a Hadamard matrix. Then

$$\|H_n\|_{\infty,2} = \max_{x \in \{\pm 1\}^n} \|H_n x\|_2 \leq \sigma_{\max}(H_n)\sqrt{n} = n.$$

So $\|Y\|_{\infty,2} \leq 1$. Let $M = n^{-1/2}I_n$. Then $\|M\|_{2,1} = 1$ by (21). Let $X \in \{\pm 1\}^{n \times n}$ be such that $\text{tr}(XH_n) = n^2$ and thus $\text{tr}(XY) = n$. Clearly $\|X\|_{1,\infty} = 1$ by (8). We have

$$\text{tr}(XMY) = n^{-1/2} \text{tr}(XY) = n^{1/2} \rightarrow \infty$$

as $n \rightarrow \infty$. □

6. CONCLUSION

We hope our characterization of Grothendieck's constant as a norm of the central object in the study of fast matrix multiplications would spur interactions between the two areas and perhaps even facilitate the determination of its exact value. Knowing that Grothendieck's inequality is a unique instance within a family of natural norm inequalities may help us better understand its ubiquity and utility. In fact, the way we formulate Grothendieck's inequality in (3) facilitated our elementary proof of the inequality in [20], which is one that works over both (i) \mathbb{R} and (ii) \mathbb{C} and yields both (iii) Krivine's bound and (iv) Haagerup's bound.

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